

Symplectic Isometries of Quantum Stabilizer Codes

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- The main theorem has connections with the LU-LC conjecture.

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- $\text{tr} : \mathbb{F}_q \rightarrow \mathbb{F}_p$ will denote the **trace** over the prime field, and $\omega = e^{2\pi i/p}$ will be a fixed p^{th} primitive root of unity.

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- The **minimum distance** of a stabilizer code is

$$\text{dist}(C) := \begin{cases} \min\{\text{wt}_s(a, b) \mid (a, b) \in C^\perp - C\} & \text{if } C \subsetneq C^\perp \\ \min\{\text{wt}_s(a, b) \mid (a, b) \in C^\perp - \{0\}\} & \text{if } C = C^\perp \end{cases}.$$

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$$\text{wt}_s(x) = \text{wt}_s(f(x)) \text{ and } \langle x | y \rangle_s = \langle f(x) | f(y) \rangle_s.$$

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- 2 $(a, b) \mapsto (\dots, a_{i-1}, b_i, a_{i+1}, \dots, \dots, b_{i-1}, -a_i, b_{i+1}, \dots)$.

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- For a linear map $f : \mathbb{F}_q^{2n} \rightarrow \mathbb{F}_q^{2n}$, denote $\tilde{f} := \gamma \circ f \circ \gamma^{-1}$.

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Theorem (Gluesing-Luerssen/P, 2017)

A linear map $f : \mathbb{F}_q^{2n} \rightarrow \mathbb{F}_q^{2n}$ is a symplectic isometry iff the map $\tilde{f} : (\mathbb{F}_q^2)^n \rightarrow (\mathbb{F}_q^2)^n$ is given by

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- We are interested on stabilizer codes.



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Open Problem

How different can the groups $\text{Mon}_{\text{SL}}(C)$ and $\text{Symp}(C)$ be?



Theorem (P, 2018)

Let $H \leq G$ be two subgroups **that satisfy some necessary conditions**. Then there exists a stabilizer code C such that $H = \text{Mon}_{\text{SL}}(C)$ and $G = \text{Symp}(C)$.

Thank You!